

# **Game reductions and FAR for infinite games**

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# Motivation I

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- Parity games with finite number of colours are positionally determined
- Muller games with finite number of colours are not positionally determined
- Finitely colored Muller games can be reduced to parity games using *latest appearance record* (LAR) or *Zielonka tree* and associated memory

# Motivation II

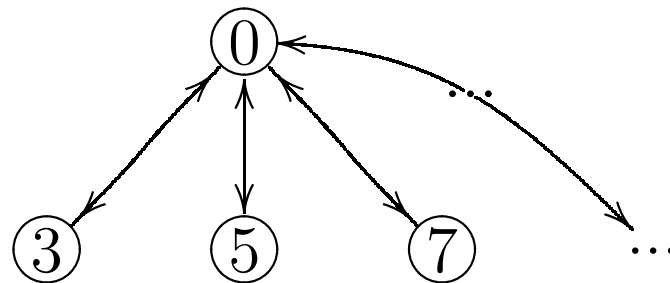
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- Infinitely ( $\omega$ ) coloured min-parity games are positionally determined
- Moreover, for every  $\omega$ -coloured positionally determined Muller game there is an isomorphic min-parity winning condition
- $\omega$ -coloured max-parity games are not positionally determined
- There are simple examples of  $\omega$ -coloured Muller games that do not have a finite memory strategy.

# Goals

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- Consider the following max-parity game:



- 0 wins with an intuitively simple strategy
- No finite memory strategy
- Goals:
  - Find a way to analyse different kinds of memory in games
  - Define a suitable memory for infinite Muller games

# Game Reductions with Memory

# Adding Memory to Games

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- *Memory* is given by a set of states  $M$  with an selected initial state  $m_0$  and an update function

$$\mu : V \times M \rightarrow M$$

- We will extend the game arena

$$G = (V_0 \sqcup V_1, E) \rightsquigarrow (V_0 \times M \sqcup V_1 \times M, E') = G'$$

- Moves

$$v \rightarrow w \rightsquigarrow (v, m) \rightarrow (w, \mu(v, m))$$

# Reducing Games

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- Extending a play

$$\pi = v_0 \rightarrow v_1 \rightarrow \dots \rightsquigarrow$$

$$\rightsquigarrow (v_0, m_0) \rightarrow (v_1, \mu(v_0, m_0)) \rightarrow \dots = \pi'$$

- $G$  reduces with memory  $\mathcal{M} = (M, m_0, \mu)$  to  $G'$  if there is a winning condition for  $G'$  such that

$$0 \text{ wins } \pi \text{ in } G \iff 0 \text{ wins } \pi' \text{ in } G'$$

- $G$  is an  $\mathcal{M}$  strategy game if it reduces with  $\mathcal{M}$  to a positionally determined game

# Using Game Reductions

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- *Example:* LAR reduction of Muller games to parity games
- Advantage: we capture different kinds of memories and prove only equivalences of plays
- Drawback: we capture only reductions for both players
- For example 0-positionality of Rabin games can not be proved directly with reductions

# Finite Appearance Record

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We consider  $\omega$ -coloured Muller games.

$(M, m_0, \mu)$  is a *finite appearance record* if

$$M = (\omega \sqcup N)^d$$

for some finite set  $N$  and number  $d$  (FAR dimension) and when

$$\mu(v, c_1, \dots, c_d) = (k_1, \dots, k_d)$$

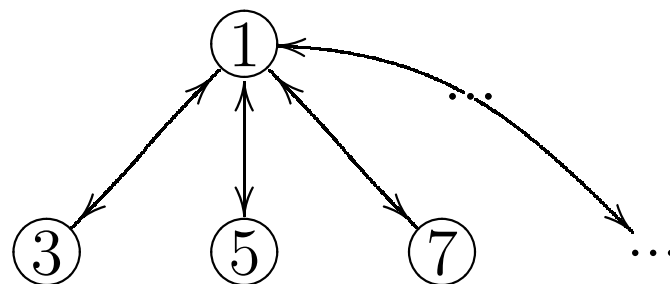
then  $k_i$  is either the colour of  $v$  or  $k_i \in N$  or  $k_i = c_j$ .

*Example:* LAR is a FAR if we take  $d$  to be equal to the number of colours in the finitely coloured game and  $N = \{\perp, B\}$ .

# FAR Properties

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- Resembles register machines, automata with infinite signature
- We can always increase the colour we see with 1-dimensional FAR
- We can infinitely often visit infinitely many colours with 2-dimensional FAR



# Reducing Infinite Games

# $\omega$ -coloured Muller Games

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- We consider infinitely coloured Muller games but with finite winning condition:

$$\mathcal{F} = \{W_1, W_2, \dots, W_N\}$$

- $W_i \subseteq \omega$  possibly infinite
- If  $\mathcal{F} = \{W_1, W_2, \dots\}$  contains infinitely many sets then we can construct a game that does not reduce to a positionally determined one with any FAR
- Since we consider only finitely many  $W_i$ 's we do not capture parity games

# Muller Games Reduction

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- *Theorem:* Muller games with finite winning condition can be reduced to finitely coloured Muller games using FAR
- For  $\mathcal{F} = \{W_1, W_2, \dots, W_N\}$  we are using FAR of dimension  $3 \cdot N$
- Reduction idea: for each  $W_i$  we use 3 registers, two to visit all colours in  $W_i$  infinitely often and the remaining one to make sure no other colours are seen repeatedly
- *Corrolary:* Each Muller game with a finite winning condition is an FAR strategy game

# Conclusions and Open Problems

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- Game reductions are a convenient method to analyse what type of memory is needed to win specific games
- Finite appearance record is enough to win  $\omega$ -coloured Muller games with finite winning conditions
- For  $\omega$ -coloured Muller games that have a Zielonka tree we can keep in memory the position on a branch of the tree and win with such memory
- Some Muller games with infinite winning condition can be won with FAR - which ones?
- Can max-parity games be won with FAR?

Thank you!